Modelling and Validation of Concurrent System

António Ravara May 6, 2024

What drives my research?

How can we (coders) make programs go right?

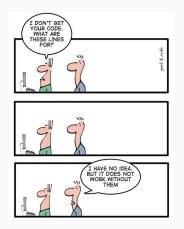
What drives my research?

How can we (coders) make programs go right? Let's take steps to avoid horror stories

What drives my research?

How can we (coders) make programs go right?

Let's take steps to avoid horror stories





Modern Software Systems

• Pervasive – our society fundamentally depends on them.

Modern Software Systems

- Pervasive our society fundamentally depends on them.
- Crucial ones are huge:
 airspace, banking, taxes, telecoms, ...

Modern Software Systems

- Pervasive our society fundamentally depends on them.
- Crucial ones are huge:
 airspace, banking, taxes, telecoms, ...
- Main characteristics: large-scale, distributed, communication-intensive, hold critical data.

Modern Software Systems

- Pervasive our society fundamentally depends on them.
- Crucial ones are huge:
 airspace, banking, taxes, telecoms, ...
- Main characteristics: large-scale, distributed, communication-intensive, hold critical data.

In short:

no room to failures!

Difficult to build and to maintain.
 Blueprints? Development methodology? Updates?

- Difficult to build and to maintain.
 Blueprints? Development methodology? Updates?
- How to express their goals and intended functionality?
 Requirements? Specifications?

- Difficult to build and to maintain.
 Blueprints? Development methodology? Updates?
- How to express their goals and intended functionality?
 Requirements? Specifications?
- How to ensure them correct?
 Do they do what they are supposed to? Will they crash?

- Difficult to build and to maintain.
 Blueprints? Development methodology? Updates?
- How to express their goals and intended functionality?
 Requirements? Specifications?
- How to ensure them correct?
 Do they do what they are supposed to? Will they crash?
- How to keep them safe? Are they hackable?

Modelling and validating (large) software systems

Modelling challenges Informal lists of requirements in natural language are not enough.

Modelling and validating (large) software systems

Modelling challenges Informal lists of requirements in natural language are not enough.

Validation challenges Testing/Debugging is not enough.

Modelling and validating (large) software systems

Modelling challenges Informal lists of requirements in natural language are not enough.

Validation challenges Testing/Debugging is not enough.

Viable approach Mathematical tools:

- represent rigorously the intended behaviour
- allow to formally verify correctness.



Dijkstra –'72 Turing award: "The humble programmer"

"Program testing can be used to show the presence of bugs, but never to show their absence!"

One striking example

Joshua Bloch, Google Research Blog (2006): "Nearly All Binary Searches and Mergesorts are Broken"

One striking example

Joshua Bloch, Google Research Blog (2006): "Nearly All Binary Searches and Mergesorts are Broken"

- A 9 years old bug on binary search in the standard Java Library
- A clear presentation on how to implement:
 Jon Bentley Programming Pearls. 1986 (2nd ed. 2000)
 The challenge of binary search

One striking example

Joshua Bloch, Google Research Blog (2006): "Nearly All Binary Searches and Mergesorts are Broken"

- A 9 years old bug on binary search in the standard Java Library
- A clear presentation on how to implement:
 Jon Bentley Programming Pearls. 1986 (2nd ed. 2000)
 The challenge of binary search

How to avoid these kind of problems:

```
Hacker-Proof Coding:
https://www.wired.com/2016/09/
computer-scientists-close-perfect-hack-proof-code/
https://cacm.acm.org/magazines/2017/8/
219596-hacker-proof-coding/fulltext
```

Pentium 5 bug (1994): rounding error

Correct value:

$$\tfrac{4,195,835}{3,145,727} = 1.333820449136241002$$

Value returned by a faulty Pentium processor:

$$\frac{4,195,835}{3,145,727} = 1.333739068902037589$$

https://en.wikipedia.org/wiki/Pentium_FDIV_bug

Success stories

Hardware verification at Intel

```
https://www.quora.com/
```

 ${\tt What-does-a-SOC-verification-engineer-at-Intel-or-AMD-accessed and the access of the access of$

Success stories

Hardware verification at Intel
 https://www.quora.com/
 What-does-a-SOC-verification-engineer-at-Intel-or-AMD-a

Driver verification at Microsoft
 https://msdn.microsoft.com/en-us/library/windows/hardware/ff552806(v=vs.85).aspx

Success stories

- Hardware verification at Intel
 https://www.quora.com/
 What-does-a-SOC-verification-engineer-at-Intel-or-AMD-a
- Driver verification at Microsoft
 https://msdn.microsoft.com/en-us/library/windows/hardware/ff552806(v=vs.85).aspx
- Specification logic (TLA) at Amazon
 http://cacm.acm.org/magazines/2015/4/
 184701-how-amazon-web-services-uses-formal-methods/fulltext

Success stories

- Hardware verification at Intel
 https://www.quora.com/
 What-does-a-SOC-verification-engineer-at-Intel-or-AMD-a
- Driver verification at Microsoft
 https://msdn.microsoft.com/en-us/library/windows/hardware/ff552806(v=vs.85).aspx
- Specification logic (TLA) at Amazon
 http://cacm.acm.org/magazines/2015/4/
 184701-how-amazon-web-services-uses-formal-methods/fulltext
- Code verification at Facebook http://fbinfer.com

Today: concurrent reactive systems

- requirements to model concurrent reactive systems
- the calculus of communicating systems (CCS)

Today: concurrent reactive systems

- requirements to model concurrent reactive systems
- the calculus of communicating systems (CCS)

Tomorrow: equivalences for CCS

- requirements for an equivalence notion
- observational behavioural equivalence

Today: concurrent reactive systems

- requirements to model concurrent reactive systems
- the calculus of communicating systems (CCS)

Tomorrow: equivalences for CCS

- requirements for an equivalence notion
- observational behavioural equivalence

Wednesday: dynamic communication topologies

- the pi-calculus (syntax and operational semantics)
- observational behavioural equivalences

Today: concurrent reactive systems

- requirements to model concurrent reactive systems
- the calculus of communicating systems (CCS)

Tomorrow: equivalences for CCS

- requirements for an equivalence notion
- observational behavioural equivalence

Wednesday: dynamic communication topologies

- the pi-calculus (syntax and operational semantics)
- observational behavioural equivalences

Thursday: logics-theory and tools

Today: concurrent reactive systems

- requirements to model concurrent reactive systems
- the calculus of communicating systems (CCS)

Tomorrow: equivalences for CCS

- requirements for an equivalence notion
- observational behavioural equivalence

Wednesday: dynamic communication topologies

- the pi-calculus (syntax and operational semantics)
- observational behavioural equivalences

Thursday: logics-theory and tools

Friday (optional): Research problems

Bibliography and resources

- R. Milner: Communication and concurrency.
 Prentice Hall 1989
- R. Milner: Communicating and mobile systems the Pi-calculus. Cambridge University Press 1999
- L. Aceto, A. Ingólfsdóttir, K. Larsen, J. Srba: Reactive systems: modelling, specification and verification.
 Cambridge University Press 2007
- C. Stirling: Modal and Temporal Properties of Processes.
 Texts in Computer Science, Springer 2001
- J. C. Bradfield, C. Stirling: Modal mu-calculi.
 Handbook of Modal Logic 2007: 721-756

Bibliography and resources

- R. Milner: Communication and concurrency.
 Prentice Hall 1989
- R. Milner: Communicating and mobile systems the Pi-calculus. Cambridge University Press 1999
- L. Aceto, A. Ingólfsdóttir, K. Larsen, J. Srba: Reactive systems: modelling, specification and verification.
 Cambridge University Press 2007
- C. Stirling: Modal and Temporal Properties of Processes.
 Texts in Computer Science, Springer 2001
- J. C. Bradfield, C. Stirling: Modal mu-calculi.
 Handbook of Modal Logic 2007: 721-756

Slides and exercises

block