

Modelling an implementation

&

implementing a model

Thanks to the organisers

Apologies to the audience

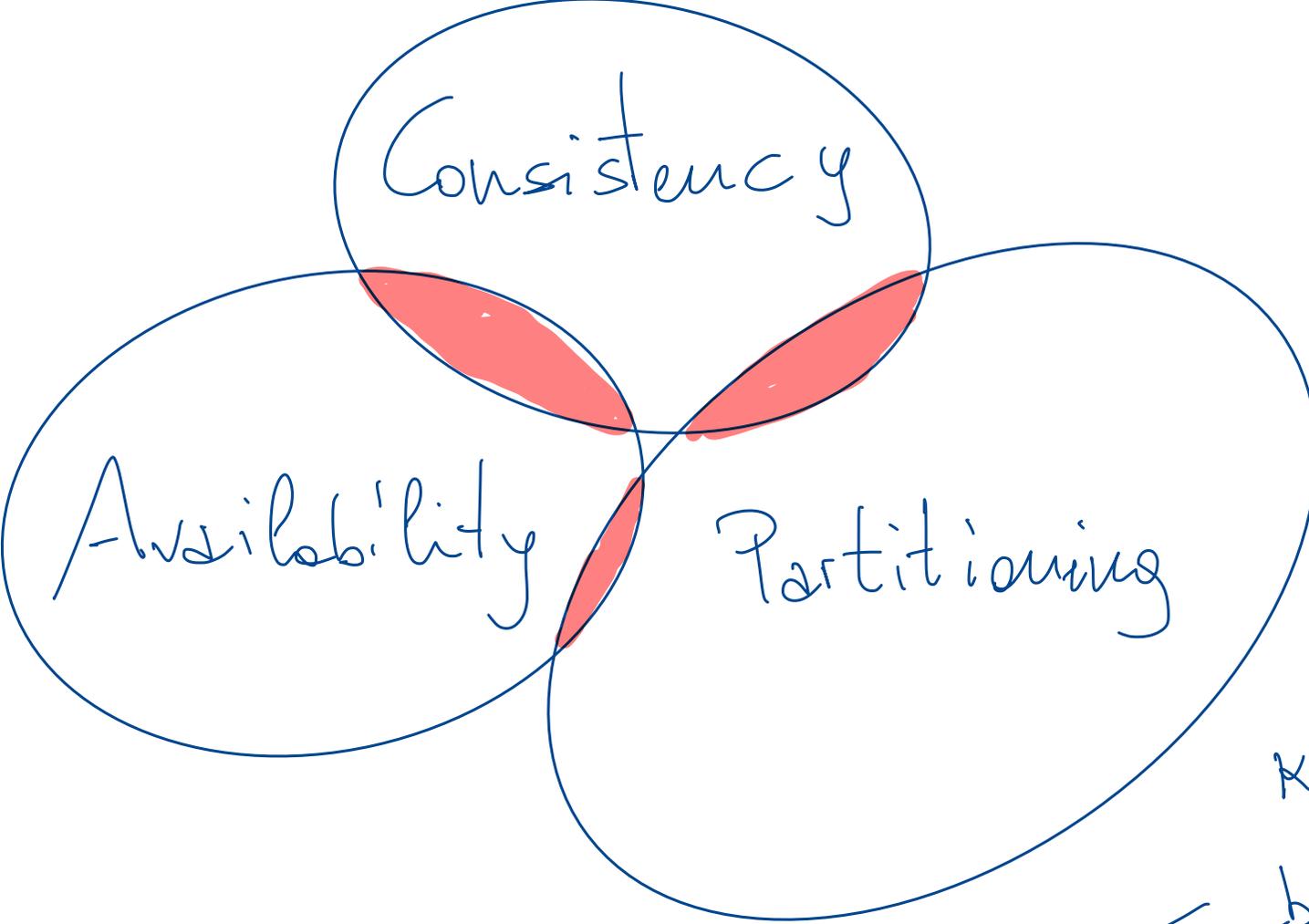
Thank Michele Pinna

Modelling and Implementation

Joint work with

Roland Kuhn @ Actyx

Hernán Melgratti @ UBA



- Systems don't stop
- Partitioning is intrinsic
- Availability is not negotiable €



Keep going and don't
be afraid to make
mistakes

LOCAL-FIRST

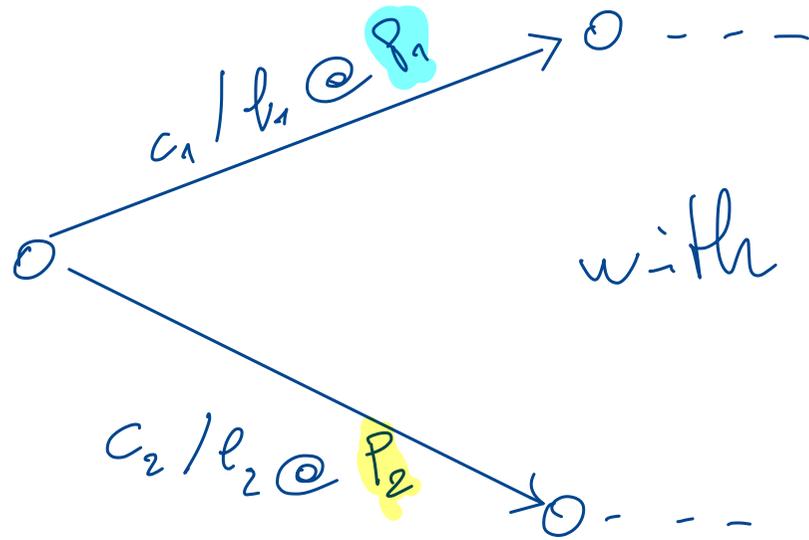
@ Actyx ~~*~~

- Platform to develop applications for controlling / managing industrial production processes
- Command execution generates events
- Pub-Sub middleware (asynchronously) replicates logs of events by merging newly generated events
- Temporary partitioning
- Need for eventual consistency of replicated logs

~~*~~ and other application domains too

Top down

- Global type (and well-formedness conditions) to handle protocols where



with $P_1 \neq P_2$ and arbitrary nom. of inst.

- Local type are FSAs too

projection op.
parametrised by
subscribers

Some comments / results ^①

- session infidelity
- Projections yield correct realisations
- Deadlock freedom by construction!
- DROPPING events is harmless in W.F. protocols

① Details in Kuhn, Melgratti, T. @ ECOOP 23)

Open Problems

- Generalise WF
- Optimise
- Support to "design" subscriptions

- Collaborative vs. Adversarial
- Failures
- Unreliable pub-sub
- impure / non-compensable events

Implementing a model

Joint work in progress with

Philip Heller @ ZTH

Herman Melgratti @ UBA

Alcriste Scelous @ DTU

Dagstuhl, September 2021

the join calculus is a powerful abstraction
for concurrent programming;

banana (b) & apple (a) if a=b \triangleright fruitSalad (b, a)
lettuce (l) & apple (a) & lives (b) \triangleright salad (l, a, b)

Q: Why hasn't the join calculus become popular?

A: Its implementations are not very efficient

Idea: reuse partial matches

Fix a signature Σ and a set of variable symbols X

μ range on $\text{Term}_{\Sigma, X}$

δ range on propositions on X

$Q \in \text{Term}_{\Sigma}^*$

pattern construction
some boolean exp with vars
queues of values

Pattern matching

$\pi = \mu_1 \& \dots \& \mu_n$ if δ

conditional
pattern

$Q \vDash_{\sigma} \pi$ where

$Q \in \text{Term}_{\Sigma}^n$

$\forall 1 \leq i \leq n: Q[i] = \pi[i]$

$\delta \sigma$

Chasing matches

Given $Q \in \text{Term}_\Sigma^n$ and $\pi = \mu_1 \& \dots \& \mu_m$

candidate
indexes

$$\text{cidx}(Q, \mu) = \{ i \mid \exists j, \sigma: Q[i] = \mu[j] \sigma \}$$

assignments

$$A = \{ \alpha \in \{1, \dots, m\}^I \mid I \subseteq \text{cidx}(Q, \mu) \text{ and } \alpha: i \mapsto j \Rightarrow \exists \sigma: Q[i] = \mu[\alpha(i)] \sigma \}$$

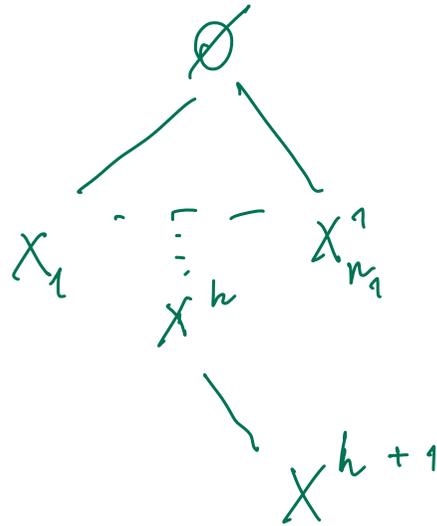
$$\alpha \subseteq \beta \Leftrightarrow \min(\text{dom } \alpha \setminus \text{dom } \beta) \leq \min(\text{dom } \beta \setminus \text{dom } \alpha) \quad \text{or } \text{dom } \alpha = \text{dom } \beta$$

Find $\alpha \in A$ "minimal" s.t.

$$\exists \sigma: \forall i \in \text{dom } \alpha: Q[i] = \mu_{\alpha(i)} \sigma \quad \text{and} \quad \gamma \sigma \text{ holds}$$

$1, \dots, r$

$r \in U$
 \vdots
 $r \in U$



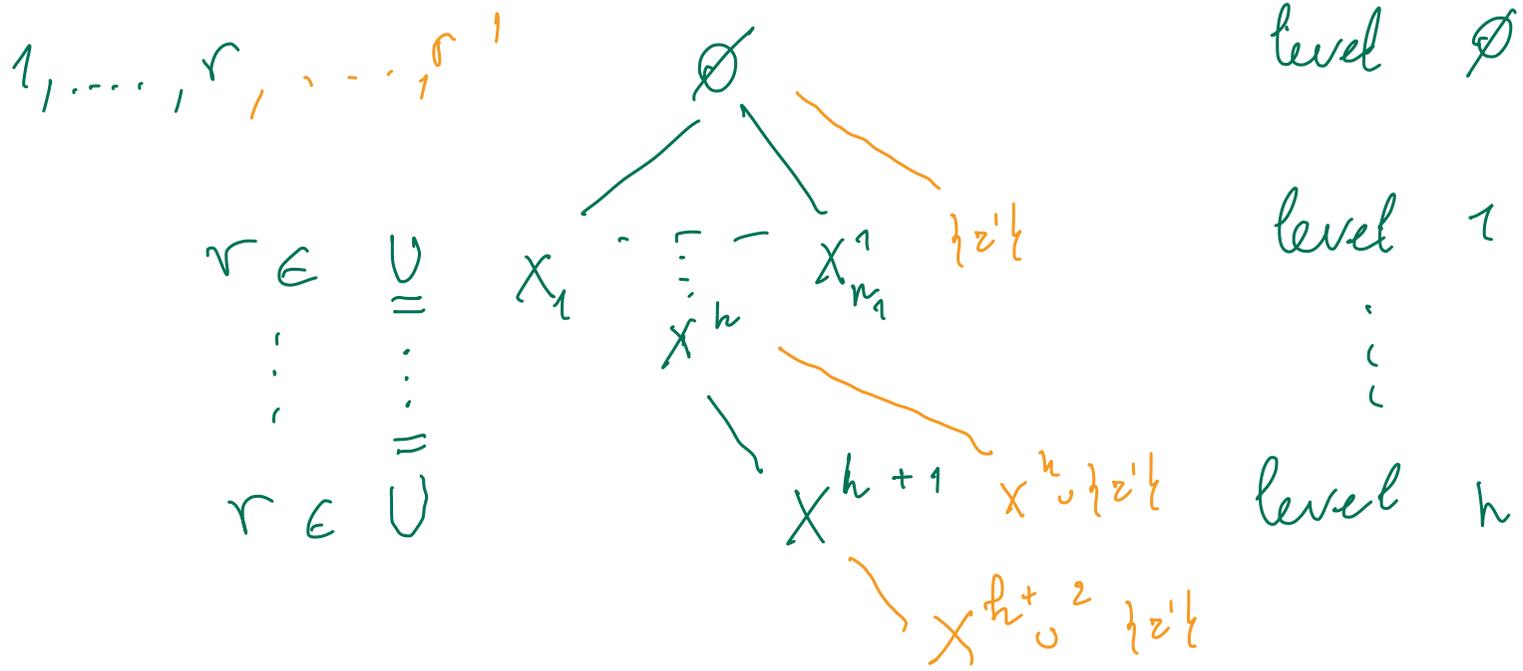
level \emptyset

level 1

\vdots

level h

X father $Y \Rightarrow \min X \leq \min Y$
 $\max X \neq \max Y$



$X \text{ father } Y \implies \min X \leq \min Y$
 $\max X \neq \max Y$

Open Problems

- Does this work better?
 - Experiments
 - Benchmarks (e.g. SAVIND)
- the algorithm is exponential in the worst case, but how realistic pathological cases are?
- Are there other interesting kinds of join patterns?
- Other notions of "best" match?

